

Arithmetic Operations on Digital and Binary Numbers (3)

Lecture 04
Book Chapter(s): 1

COE211-Digital Logic Design

الفصل الدراسي الأول 1442-2020 Fall

جامعة طيبة فرع ينبع - كلية علوم وهندسة الحاسبات - شطر الطالبات



جامعة طيبة

د. فاطمة الحربي

1s Complements - Addition and Subtraction (1)

- **Algorithm for addition, $A + B$:**

1. Perform binary addition on the two numbers
2. If there is a carry out of the MSB, add 1 to the result
3. Check for overflow:
 - Overflow occurs if result is opposite sign of A and B

- **Algorithm for subtraction, $A - B$:**

1. $A - B = A + (-B)$
2. Take 1s complement of B by inverting all the bits
3. Add the 1s complement of B to A

1s Complements - Addition and Subtraction (2)

- Examples: 4-bit binary system

$$\begin{array}{r} +3 \quad \quad 0011 \\ + +4 \quad + 0100 \\ \hline +7 \quad \quad 0111 \\ \hline \end{array}$$

$$\begin{array}{r} +5 \quad \quad 0101 \\ + -5 \quad + 1010 \\ \hline -0 \quad \quad 1111 \\ \hline \end{array}$$

$$\begin{array}{r} -2 \quad \quad 1101 \\ + -5 \quad + 1010 \\ \hline -7 \quad \quad \color{red}10111 \\ \hline \quad \quad + \quad 1 \\ \hline \quad \quad 1000 \end{array}$$

$$\begin{array}{r} -3 \quad \quad 1100 \\ + -7 \quad + 1000 \\ \hline -10 \quad \color{red}10100 \\ \hline \quad \quad + \quad 1 \\ \hline \quad \quad 0101 \end{array}$$

2s Complements - Addition and Subtraction (1)

- **Algorithm for addition, $A + B$:**

1. Perform binary addition on the two numbers
2. Ignore the carry out of the MSB (most significant bit)
3. Check for overflow:
 - Overflow occurs if the 'carry in' and 'carry out' of the MSB are different, or if result is opposite sign of A and B

- **Algorithm for subtraction, $A - B$:**

1. $A - B = A + (-B)$
2. Take 2s complement of B by inverting all the bits and adding 1
3. Add the 2s complement of B to A

2s Complements - Addition and Subtraction (2)

- Examples: 4-bit binary system

+3	0011
+ +4	+ 0100
----	-----
+7	0111
----	-----

-2	1110
+ -6	+ 1010
----	-----
-8	1 1000
----	-----

+6	0110
+ -3	+ 1101
----	-----
+3	1 0011
----	-----

+4	0100
+ -7	+ 1001
----	-----
-3	1101
----	-----

Overflow (1)

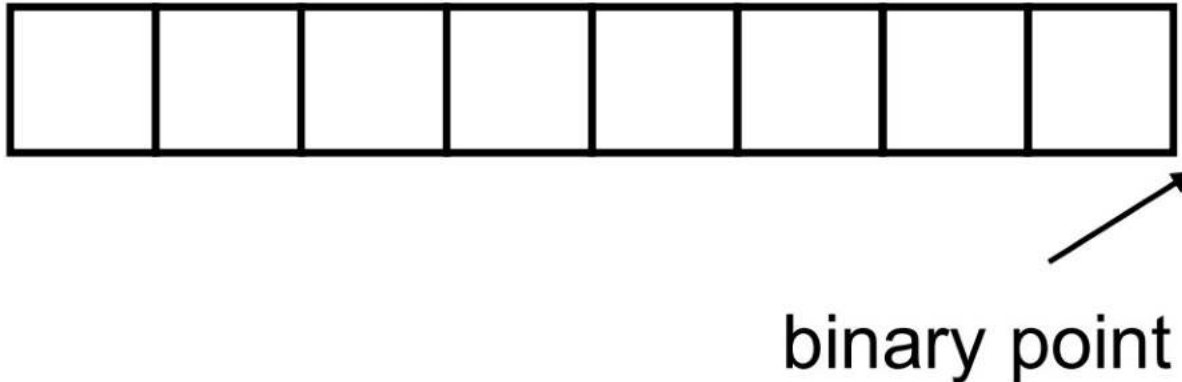
- Signed binary numbers are of a fixed range
- If the result of addition/subtraction goes beyond this range, overflow occurs
- Two conditions under which overflow can occur are:
 - *positive add positive gives negative*
 - *negative add negative gives positive*

Overflow (2)

- Examples: 4-bit numbers (in 2s complement)
- Range : $(1000)_{2s}$ to $(0111)_{2s}$ or $(-8)_{10}$ to $(7)_{10}$
 - $(0101)_{2s} + (0110)_{2s} = (1011)_{2s}$
 $(5)_{10} + (6)_{10} = -(5)_{10}$?! (overflow!)
 - $(1001)_{2s} + (1101)_{2s} = (10110)_{2s}$ discard end-carry
 $= (0110)_{2s}$
 $(-7)_{10} + (-3)_{10} = (6)_{10}$?! (overflow!)

Fixed Point Numbers (1)

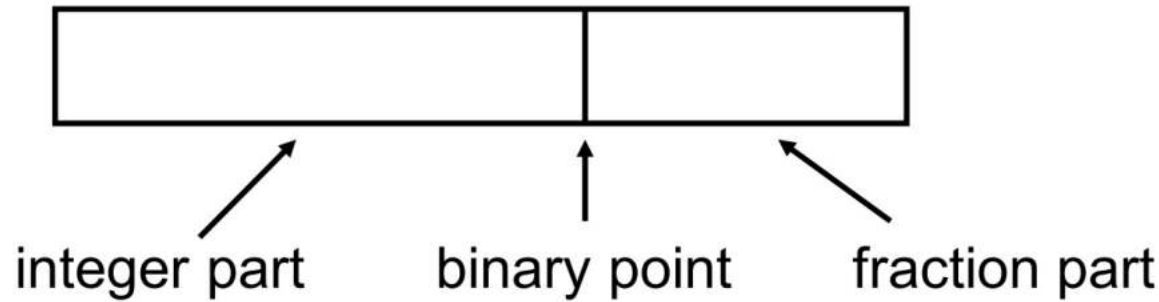
- The signed and unsigned numbers representation given are **fixed point numbers**
- The **binary point** is assumed to be at a fixed location, say, at the end of the number:



- Can represent all integers between -128 to 127 (for 8 bits).

Fixed Point Numbers (2)

- In general, other locations for binary points possible



- Examples: If two fractional bits are used, we can represent:

$$(001010.11)_{2s} = (10.75)_{10}$$

$$(111110.11)_{2s} = -(000001.01)_2$$
$$= -(1.25)_{10}$$

Floating Point Numbers (1)

- Fixed point numbers have limited range
- To represent very large or very small numbers, we use floating point numbers (cf. scientific numbers).
- Examples:

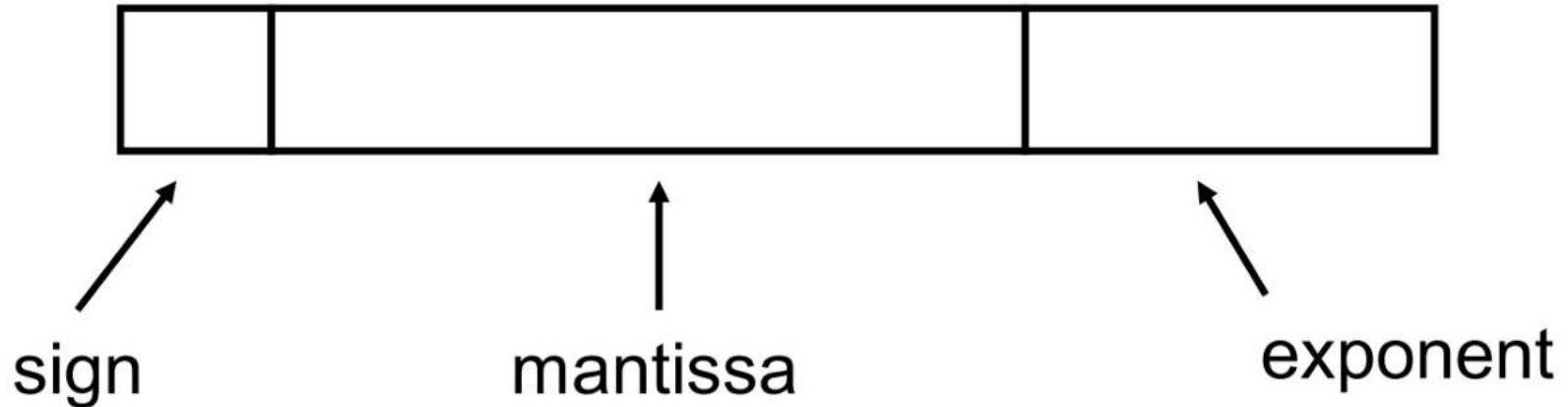
0.23×10^{23} (very large positive number)

0.5×10^{-32} (very small positive number)

-0.1239×10^{-18} (very small negative number)

Floating Point Numbers (2)

- Floating point numbers have three parts:
sign, **mantissa**, and **exponent**
- The base (radix) is assumed (usually base 2)
- The sign is a single bit (0 for positive number, 1 for negative).



Floating Point Numbers (3)

- Mantissa is usually in normalised form:
 - (base 10) 23×10^{21} normalised to 0.23×10^{23}
 - (base 10) -0.0017×10^{21} normalised to -0.17×10^{19}
 - (base 2) 0.01101×2^3 normalised to 0.1101×2^2
- **Normalised form:** The fraction portion cannot begin with zero
- More bits in exponent gives larger range
- More bits for mantissa gives better precision

Floating Point Numbers (4)

- Exponent is usually expressed in complement or excess form
- Example:
 - Express $-(6.5)_{10}$ in base-2 normalised form:
 $-(6.5)_{10} = -(110.101)_2 = -0.110101 \times 2^3$
- Assuming that the floating-point representation contains 1-bit sign, 5-bit normalised mantissa, and 4-bit exponent.
 - The above example will be represented as:

1	110101	011
---	--------	-----

Arithmetic with Floating Point Numbers - Multiplication

- Arithmetic is more difficult for floating point numbers
- Steps for multiplication:
 1. Multiply the mantissa
 2. Add-up the exponents
 3. Normalise

1. Example:

$$\begin{aligned} & (0.12 \times 10^2)_{10} \times (0.2 \times 10^{30})_{10} \\ &= (0.12 \times 0.2)_{10} \times 10^{2+30} \\ &= (0.024)_{10} \times 10^{32} \quad (\text{normalise}) \\ &= (0.24 \times 10^{31})_{10} \end{aligned}$$

Arithmetic with Floating Point Numbers - Addition

- Steps for addition:
 1. Equalise the exponents
 2. Add-up the mantissa
 3. Normalise

- Example:

$$\begin{aligned} & (0.12 \times 10^3)_{10} + (0.2 \times 10^2)_{10} \\ &= (0.12 \times 10^3)_{10} + (0.02 \times 10^3)_{10} \quad (\text{equalise exponents}) \\ &= (0.12 + 0.02)_{10} \times 10^3 \quad (\text{add mantissa}) \\ &= (0.14 \times 10^3)_{10} \end{aligned}$$

Can you figure out how to perform
SUBTRACTION and **DIVISION** for
(binary/decimal)
floating-point numbers?